

Assembly - 8.28.11

CSCI 2160

* Hollerith Code

- 1880 census took 10 years to calculate
- Hollerith came up tabulating mach
- the card had 80 columns
- 12-1 was A, 12-2 was B
- I was 12-9
- only did uppercase
- J was 11-1, R was 11-9
- S was 0-2, Z was 0-9

* Number Systems

$$- 476_9 + 583_9$$

$$\begin{array}{r} 476_9 \\ + 583_9 \\ \hline 1110_9 \end{array}$$

$$467_8 + 557_8$$

$$\begin{array}{r} 467_8 \\ + 557_8 \\ \hline 1246_8 \end{array}$$

$$\begin{array}{r} 1011_2 \\ + 1101_2 \\ \hline 11000_2 \end{array}$$

$$\begin{array}{r} 312_{15} \\ 888_9 \\ - 146_9 \\ \hline 256_9 \end{array}$$

$$\begin{array}{r} 1011_{11} \\ 214_4 \\ - 132_4 \\ \hline 13_4 \end{array}$$

$$!_{10}H = 2560 + 0 + 11 = 2571_{10}$$

- needed 8 bits in a byte so that all characters on a keyboard could be represented
- addresses are binary also
- $1100_2 = 12$

* EBCDIC

- 12-1 - $1100_2 | 0001_2$ - A
- 12-2 - $1100_2 | 0010_2$ - B

- when reading memory, you need to know context

- 11010001 - J in EBCDIC
- 11100001 - S in EBCDIC
- EBCDIC has gone and digit, derived from Hollerith

* ASCII

- starting with capital A, 0100 0001 (65 in decimal) + 4_h
- letters after are incremented
- for ASCII, turn on 3rd bit
- lowercase a = 0110 0001 - 97 + 6_h
- for lowercase, add 32 in decimal or 20_h

* Numbers in EBCDIC

- 0 - F0 in EBCDIC
- 9 - F9 in EBCDIC
- EBCDIC, characters precede numbers when sorting
- ASCII - numbers precede characters

* Numbers in ASCII

- 0 is 30_h in ASCII
- 9 is 39_h in ASCII

CSCI 2160

- byte = 8 bits 2^8

- word = 2 bytes 2^{16}

- double word = 2 words = 4 bytes 2^{30}

- quad word = 4 words = 8 bytes 2^6

- MSB is the sign bit

- use two's complement - first represent positive value, flip the bits & add 1

$$-37 = \emptyset \text{ DBh}$$

- 97 $\rightarrow$ 01100001 $\rightarrow$ 1001110 + 1 = 1001111 = 9Fh

$$- \quad - \quad | \rightarrow 00000001 \rightarrow 11111110 + 1 = 11111111 = \text{FFFF}$$

- in Assembly, you must write a 0 before a character in hex

- 07h as unsigned = 215

- D7h no character = 6

$$D7h \text{ as signed} = 28h \text{ is flipped} = 28h + 1h = 29h = -41$$

- 83h as resigned = $128 + 3 = 131$

$$-83h \text{ as signed} = 7Ch + 1 = 7Dh = -125$$

- 73h as unsigned = $112 + 3 = 115$

$$- 734 \text{ as signed} = 115$$

- 80h as unsigned = 128

- 80h as signed = -128

- you can always represent one more negative value
the positive is a signed binary number

- 8000h as signed = -32768

- 8000h as unsigned = 32768

- 125 as byte = $01111101 + 1 = 01111110 \rightarrow 11111110 \rightarrow 0FDh$ X

-125 as byte = 01111101 $\xrightarrow{\text{bits}}$ 10000010 + 1 = 10000011 $\xrightarrow{\text{flip bits add 1}}$ 83h $\checkmark$

* Symbols

- byte - db
- word - dw
- double word - dd
- quad word - dq
- answer byte
- byte answer; Java
- int value = 35; Java
- value dd 35 Assembly
- or value dword 35 Assembly
- Java → Assembly
- short tax = 13; → tax dw 13h
- long value = 19; → value qw 19h
- int number = -20 → number dd -20
- byte answer = 'B'; → answer db 'B'

(write lined up)

.data

tax word 13

value qword 19

number dword -20

answer byte 'B' (you can use ' or ")

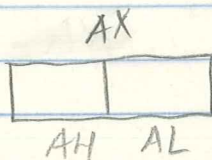
- store values in memory in reverse byte order
- Assembly assumes decimal if there is no trailing letter; use d to specify decimal, use h to specify hex
- 13d → ^{hex}000Dh → ^{reverse byte}0D00h
- 20d as dword → ^{hex}00000014h → ^{flip bits + 1}FFFFFFEh → ^{reverse byte}EFFFFFFFh
- 'B' as byte → ^{hex}42h → ^{reverse}42h

- | statement | hex | reverse byte |
|----------------|---------|--------------|
| - A word 22 → | 0016h | → 1600h |
| - B word 19 → | 0013h | → 1300h |
| - C dword 37 → | 000025h | → 25000000h |
- data values are stored sequentially
 - strings are not stored in reverse byte order

* Registers

- RAX, RBX, RCX, RDX (workhorse registers) (32-bits)
- CS, DS, SS (32 bit)
- flat segment model - can use all 32 bits to represent an address
- 32 bits can store 4 GB of addresses
- 2^{10} - terabyte
- 2^{50} - petabyte
- 2^{60} - exabyte
- 2^{70} - zettabyte
- 2^{80} - yottabyte
- lower RAX is EAX, lower RBX is EBX, so on
- descriptor table

code	data	stack
0EC2FB13	0B1CF2E8	1C2031D3
↑	↑	↑
CS	DS	SS



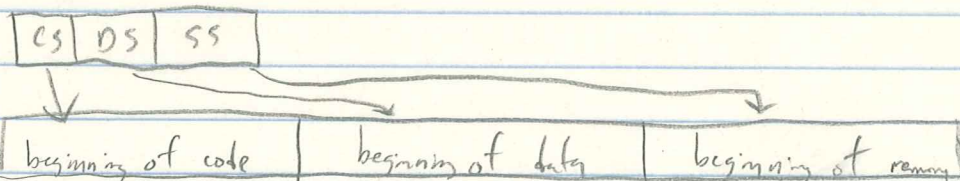
AX is 16 bits
 AL/AH is 8 bits
 EAX is 32 bits

* Quiz 1 & 2 Review

- A in ASCII is 41h, 65d
- a in ASCII is 61h, 97d
- A in Hollerith is 12-1, EBCDIC is C1h
- J in Hollerith is 11-1, EBCDIC is D1h (not direct translation)
- C1-C9, D1-D9, E2-E9 - EBCDIC alphabet
- 1d in EBCDIC is F1h, ASCII is 31h
- 0d in ASCII is 30h
- Remember that a digit in a base can not be larger or equal to the base
- Long quiz on Thursday! 9.8.11

* Notes

- register - high speed memory that perform arithmetic, holds instructions
- warehouse registers - A, B, C, & D
- 8 bit - AH & AL
- 16 bit - AX, BX, CX, & DX
- 32 bit - EAX, EBX, ECX, EDX
- 64 bit - RAX, RBX, RCX, RDX
- CS, DS, SS - 16 bit
- CS - code, DS - data, S - segment (stacks)
- CS gives us a description



- index registers - SI & DI
- SI - source index, DI - destination index
- used to hold starting address of something

Assembly

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(SCI 2162)

- index registers are 16 bit
- ESI, EDI - extended index registers, 32 bit
- RSI, RDI - 64 bit index registers
- pointer registers (no 8 bit)
- BP - base pointer, keeps track of where stack was at time a piece of code ^{is} entered
- SP - stack pointer, points to top of stack
- IP - instruction pointer
- IP is not to be messed with, but hackers uses IP to inject code
- IP is updated automatically by operating systems
- all of these have extended registers - EBP, ESP, EIP
- .data, .stack, .code - directives
- .stack 100h - allocates stack of 256 bytes

- 1 byte db
- 2 word dw
- 4 dword dd
- 8 qword dq

			location counter
bVar	byte	?	0000
wVal	word	27	0001
dVal	dword	-7	0003
qNum	qword	19	0007
			000F

- assembler creates symbol table
- symbol - location counter - size

Assembly

CSCI 2160

- if you define an array of values, the location counter reflects that 1 to 1
- the size table reflects the size of what is declared (byte, word, dword, etc.)

loc ctr	.data	equivalent hex	size
0000	byte 9, 7	09, 07	1
0002	word 17, -2, 3	0011, FFFE, 0003	2
0008	dword 29, 32	0000001D, 00000020	4
0010	qword 5	00000000 00000005	8
0018	byte '123A'	31 32 33 41	1
001C			

- in memory ...

.data

09 07

11 00 FE FF 03 00

1D 00 00 00 20 00 00 00

05 00 00 00 00 00 00 00

31 32 33 41

- dup directive (used with # in front)
- 2 dup (?) - 00 00 (byte)
- 3 dup (7) - 0007 0007 0007 (word)

Assembly

CSC 2160

- wX word 33, 19
 wY word 26, 39, 55
 wZ word ?
 wA word 2 dup (5)
 dVal dword 3 dup (-19)
 dStr byte 2 dup ('ABC')

loc chr	actual hex	reverse byte order
0000	0021, 0013	21 00 13 00
0004	001A, 0027, 0037	1A 00 27 00 37 00
000A	0000	00 00
000C	0005, 0005	05 00 05 00
0010	FFFFFFED, FFFFFFFED, FFFFFFFED	ED FF FF FF ED FF FF FF ED FF FF FF
001C	41 42 43 41 42 43	41 42 43 41 42 43
0022		

- Quiz on Thursday will cover everything up to 9.6.11

Assembly

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- ESP contains address of stack first byte beyond the top
- on Quiz 3, 12b is 0018ff8c
12c is 0018FE8C
- push - 1.) subtracts from ESP 2.) copies bytes to stack
- pop - 1.) copies data from stack 2.) pointer moved
- Quiz on Thursday, 9.15

- Integer constants

decimal - add ax 24; add ax 24d

binary - add ax 100b

hex - add ax 64h

character - add ax 'A'

string - add ax 'John went'; add ax "John went"

- Integer expressions

() - parenthesis

* & / - multiply & divide

mod - modulus

+ & - add, subtract

- Reserved words

instruction mnemonics - add, sub, mul, mov

registers - eax, ebx...

directives - tells MASM what to do

attributes - messages to assembler - byte, dword

operators - +, -, ()

pre-defined symbols - @

- Identifiers

must begin with A-Z, a-z, -, \$, @, ?

can be up to 255 characters

usually not case sensitive, unless switch is turned on

must start A-Z, a-z, other characters can be numerical

Assembly

- Directives

not executable

define variable or procedures

can assign names to memory segments or code

case insensitive

example - .486, .model

first three lines of program should be

.486

.model flat

.stack 100h

- .asm → assembler ↔ .lst
 .obj → linker
 .obj
 .lib

- Instructions - 4 parts

[label] operation [operands] [comment]

label - assigned address of the instruction that follows, label appears in the symbol table

- to debug, -start:

first line under .code

- last thing is PUBLIC start

- last line is END

operation - mov

- b for byte, w for word, d for dword, q for quad,
 str for string (use \$ for variable names)

- CRLF - 10 13, 0

- str Prompt1 byte "Enter first number", 0

- to use Bailey's macro, follow string with a 0

- receiving operand is always first

- each instruction must have a comment!
- `add ebx, 4` ; ebx points to next integer value

```
; Program
; Programmer
; Course
```

Write why? for comments

- Assembly cannot do memory to memory operations, must use registers!
- `add wX, wY` - WRONG!
- `mov ax, wX` ; wZ =
- `add ax, wY` ; wX +
- `mov wZ, ax` ; wY
- `sub op1, op2` ; op1 gets result of $op1 - op2$
- Nop

`mov eax, dVal` ; dVal is 32b

loader aligns into one it is as a double word boundary
 nop is used to align, takes up one byte of storage,
 speeds up processing

- Quiz 4 will be over this review

```
; Program      filename.asm      line up!
; Programmer   Name
; Course       (SCI 2160-201)
; Assignment
; Date
; Purpose
```

```
.486
.model flat
.stack 100h
```

(review handout on D2L)

- display prompt 1
- get value 1
- display prompt 2
- get value 2
- add values
- convert answer
- display answer

- .code

Quiz 4 BwB.asm

- start:

```
puts strPrompt1      ; display first input prompt
gets strInput, 10     ; get first value from keyboard
stringToNum strInput, dXval ; convert to a true binary value
puts strPrompt2      ; display 2nd input prompt
gets strInput, 10     ; get second value from keyboard
stringToNum, dXval    ; convert to a true binary value
mov eax, dXval        ; dXval =
add eax, dYval        ; dYval +
mov dXval, eax        ; dXval
numToString strOut, dXval ; convert the binary value to a numeric string
puts strOut          ; display string
```

```
INVOKE ExitProcess, 0
PUBLIC start
END
```

- from memory to register, remember that it is stored in memory in reverse byte order
- to reference memory, use brackets
 - mov ebx, offset wT ; address of wT
 - mov eax, [ebx] ; load address of ebx into eax
- | | | |
|----|-------|----|
| 24 | 37/AE | 57 |
|----|-------|----|

 - mov ax, wT
 - mov ax, wT + 1 AE37
 - mov ax, wT[1] - 1 reference byte, not position (offset)
- wT word 5, 4, 3, 2, 17, 22
 short wT = {5, 4, 3, 2, 17, 22}
 - mov ax, wT ax: 0005
 - mov ax, wT[1] is the same as mov ax, wT + 1 ax: 0000
 - mov ax, wT[4] ax: 0003
- mov eax, [ebx + esi] - add two registers to determine address
 must use ebx & esi - class standard
- ml /c /coff /Zi /Fl test.asm
 micro.bat
 " " %1.asm
- link /debug /subsystem:console /entry:start /out:%1.exe %1.obj
 convut.1201109.obj kernel32.lib utility201109.obj

Assembly

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- equal-sign directive

COUNT = 10

COUNT is replaced, works as a constant, compiler substitutes that value at assemble time

- \$ - current value of location counter

wVal word 10, 20, 18, -7, 55

00AC

wValSize = (\$ - wVal) / # of bytes (2)

0006

mov ecx, wValSize

(use ecx for looping)

- equ directive - equates

xVal equ [ebp + 4]

equ is used more than = or \$

- movsx - copy & extend sign

mov eax, wVal

FFFFFFFF

- movzx - copy & extend zero

mov eax, wVal

0000FFFF

wVal word -1

- cbw - convert byte to word

AL → AX

- cwd - convert word to double word

AX → DX:AX pair

- cwide - convert word to double word extended

AX → EAX

- cdq - convert double word to quad word

EAX → EDI:EAX pair

works only for ax/edx registers

- EFLAGS - 32 bits

status - contained in lower 16 bits

15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Assembly

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- 0 - carry flag
- 2 - parity flag
- 6 - zero flag
- 7 - sign flag
- 11 - overflow flag

carry - set when there is a carry out on an operation
 parity - set to keep number of 1 bits odd for odd parity machine
 sign - set to 1 when result is negative
 zero - set to 1 when result is 0
 overflow - result to 2 positives becoming negative

add ax, val

jo - jump overflow

jc - jump carry

- inc [operand] → value is increased by 1 (doesn't touch flags)

inc [ebx] → invalid! must know size

inc word ptr [ebx] - valid

- neg [operand] → replace operand with 2's complement (in flag set)

neg ax

- add al, bVal al = 127 bVal = 1

CF = 0, OF = 1

add -128 + -128

CF = 1, OF = 1

handle overflow by jo

- EFL represents the EFLAGS
EFL will turn red when it is changed
- ESP - stack pointer
- neg does not just flip bits, it performs a 2's complement

Assembly

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- Exam 1 is next Thursday
- cbw - works with A register - convert byte to word

eax: CE24E5C6

cbw = CE24FFC6

cwd - convert word to double word dx: ax

eax: CE24E5C6 edx: 2416ABCD

cd = edx = 2416FFFF

cwde - eax: FFFFE5C6

cdq - edx: eax

FFFFFFFF: CE24E5C6

movsx eax, cl

cl goes into al & sign is filled

movzx eax, cl

cl goes into al & zero is filled

Assembly

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Page 1

- type - # of bytes - the type that is assigned
- lengthof - # of different values
- sizeof - total # of bytes used

bVal Byte 20 dup(0)

type=1, length=20, sizeof=type x lengthof

```
mov ax, type bVal    ax: 00 01
mov ax, lengthof bVal  ax: 00 14
mov ax, sizeof bVal   ax = 00 14
```

- label - doesn't change location etc., gives a location a different type

dT dword ?

loc: 00 4C

- jump [label] - unconditional jump - "go to"

```
mov eax
add eax, dVal
jump end-of-program
```

- loop - works with ecx
 - 1) ecx is decremented by 1
 - 2) if (ecx != 0), then jump to the label

- wVal word 15, 20, 32, -1, 6

```
mov ecx, lengthof wVal;
mov eax, 0;
mov ebx, offset wVal;
```

ecx = 5

set sum to 0;

ebx = beginning of array;

stloop:

add eax, word ptr [ebx];

sum += element at ebx

add ebx, 2;

ebx points to next element

loop stloop;

- strOne byte "Hello, there!", 0
strTwo byte 40 dup(?)

mov ecx, Sizeof strOne;

mov eax, 0;

mov ebx, offset strOne;

mov edx, offset strTwo;

stloop:

mov al, byte ptr [ebx];

mov strTwo, al;

add ebx, type strOne;

add edx, type strTwo;

loop stloop

- use ebx, esi, & edi for pointer registers
- inc & dec is a little faster than add

```
- [label] proc
    push ebp
    move ebp, esp
```

```
    pop ebp
    ret
end proc
```

; this is the shortest proc, in Java, this is { }

```
; short function (short, short)
```

```
    wVal1 word ?
    wVal2 word ?
    wVal3 word ?
```

```
; wVal3 = function (wVal1, wVal2)
```

```
    push wVal2
```

```
    push wVal1
```

```
    call function
```

```
    mov wVal3, ax
```

```
; or mov wVal3, word ptr eax
```

```
    add esp, 4
```

```
; rVal = fl(iVal, sVal, iVal2)
```

```
    push sVal2
```

```
    push rVal1
```

```
    push iVal
```

```
    call fl
```

```
    mov rVal, ax
```

```
    add esp, 8
```

```

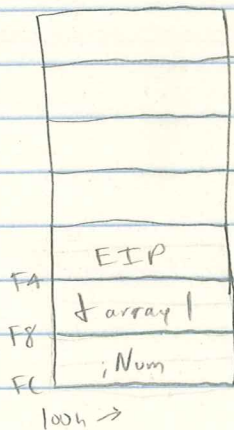
- proble SumArray(&array, iNum) ; C syntax
  push iNum
  push offset array
x  call SumArray
  mov iTotal, eax
  add esp, 8

```

- 1) EIP is pushed onto stack
- 2) EIP is loaded with address of procedure
when the function is entered, the old EIP value is on the top of the stack
- 3) when the function is entered, ebp is pushed onto stack

- the programmer must push & pop the stack equal times by the end of the method

- 4) ret copies what is on top of stack to EIP



- $sZ = \text{addshorts}(sX, sY)$

wX word 27

wY word 16

wZ word ?

; wZ = addshorts(wX, wY)

push wY

push wX

call addshorts

mov wZ, eax

add esp, 4

addshorts proc

push ebp

mov ebp, esp

mov ax, [ebp+10]

add ax, [ebp+8]

pop ebp

ret

- strCopy(dstOne, dstTwo, iSize)

push iSize

push offset dstOne

push offset dstTwo

call strCopy

add esp, 8;

; restore stack pointer

```

strCopy proc
    push ebp                ; preserve base pointer
    mov ebp, esp            ; new stack frame
    push eax; push ecx; push edi; push esi;    ; preserve registers
    mov ecx, [ebp+16]        ; ecx = # of bytes copied
    mov esi, [ebp+12]        ; esi → source str
    mov edi, [ebp+8]         ; edi → dest str

stLoop:
    mov al, [esi]
    mov [edi], al
    inc esi
    inc edi
    loop stLoop

```

EC	EBP	[ebp]
ED	EIP	[ebp+4]
EB	dstTwo	[ebp+8]
ED	dstOne	[ebp+12]
EC	:Size	[ebp+16]

- int AddUp (t arrayints, iNum) ; write procedure

```

AddUp proc
    ✓ push ebp
    ✓ mov ebp, esp
    ✓ mov ecx, [ebp+12]
    mov esi, [ebp+8] ; mov eax, 0;

```

[ebp]	EBP
[ebp+4]	EIP
[ebp+8]	t arrayints
[ebp+12]	iNum

```

stLoop:
    mov ebx, [esi] ; move source into ebx
    add eax, ebx   ; add eax to ebx
    add esi, type arrayints ; increment esi by # of bytes in type
    loop stLoop    ; restart loop if ecx ≠ 0

```

would also want to preserve ebx & ecx after AddUp proc

- Compare

cmp op1, op2

what happens is $op1 - op2$ so that flags get set
still cannot compare memory to memory

strName byte "Hello, Here!", 0

mov ebx, offset strName ; ebx $\rightarrow$ array

cmp byte ptr [ebx], 0

$\left(\begin{array}{l} \text{je} = \text{jump equal} \\ \text{jne} = \text{jump not equal} \\ \text{jl} = \text{jump less than} \end{array} \right)$ je tests the zero flag
 je finished

- sub count, 1

jz done

jz tests the zero flag

je = jz but use which one makes sense

- msgHello db " — ", 0
puts

loopStr:

mov ecx, 100 ; counter
x CRLF db 10, 13, 0 ✓ puts CRLF
x msgHello db " — ", 0 ✓ puts msgHello
~~puts msgHello~~
loop loopStr

- f1(int, short, int) : int UML
iVal = f1(iVal1, sVal2, iVal3) Java

push iVal3
push sVal2
push iVal1
call f1
mov iVal, eax
add esp, 10

- sVal = f1(sVal1, sVal2, bVal1) ; must cast

mov al, bVal1
cbw ; ax = bVal1
cwd ; eax = bVal1
push eax
push sVal2
mov ax, sVal1
cld
push eax

call f1
mov sVal, eax
add esp, 10

- $f1(int1, short, int) : int$ UML
 $iVal = f1(\&array1, sVal, \&array2)$ Java

push offset array2
 push sVal
 push offset array1
 call f1
 mov iVal, eax
 add esp, 10

EE	EBP	[ebp]
FJ	EIP	[ebp+4]
F6	int1	[ebp+8]
FA	short	[ebp+12]
FC	int2	[ebp+14]

- $f2(int, short) : void$
 $f2(iVal, sVal);$

push sVal
 push iVal
 call f2
 add esp, 6

Fa	EBP	[ebp]
F6	EIP	[ebp+4]
FA	iVal	[ebp+8]
FE	sVal	[ebp+12]

- $f2(int, short) : void$
 $f2(sVal, bVal)$

mov al, bVal
 cbw ; ax = bVal
 push ax
 mov ax, sVal
 cwd ; eax = sVal
 push eax
 call f2
 add esp, 6

- mul - multiply two positive (even if one is negative)
- imul - sign is considered

imul op

can only multiply 1x1, 2x2, or 4x4

iVal = bVal x sVal ; cannot do!

what kind of multiply is determined by operand

op is multiplies

mul bVal 1 byte multiplier → al, result is in ax

imul sVal 2 byte ... → ax, result is dx:ax

imul dVal 4 byte → eax, result is edx:eax

bVal byte 5

sVal word 10

iVal dword 15

sVal/2 = bVal/2 x bVal/2 ; legal

iVal = sVal/1 x sVal/2 ; legal

dVal = iVal/1 x iVal ; legal

- 4 x 5

mov al, 4

mov bl, 5

imul bl

; ax = 20

- 300 x 10

mov ax, 300

mov bx, 10

imul bx

; dx:ax = 3000

- $iVal = sVal1 \times sVal2$

```
mov ax, sVal1
imul sVal2      ; dx:ax
mov word ptr iVal, ax
mov word ptr iVal+2, dx
```

or

```
mov ebx, offset iVal
mov [ebx], ax
mov [ebx+2], dx
```

- $sVal = sVal1 \times sVal2$

```
mov ax, sVal1
imul sVal2      ; dx:ax
mov sVal, ax    ; dx is lost
```

- $LVal = iVal1 \times iVal2$

```
mov eax, iVal1
imul iVal2      ; edx:eax
mov dword ptr LVal, eax
mov dword ptr LVal+4, edx
```

or

```
mov ebx, offset LVal
mov ebx, eax
mov ebx+4, edx
```

- byte x byte → anything but sign overflows into AH, (CF + OF is set)
- word x word → dx, (CF + OF is set)
- dword x dword → edx, (CF + OF is set)

- iVal = sVal x bVal ; convert byte to short

```
mov al, bVal
cbw      ; ax = bVal
imul sVal ; dx:ax
mov ebx, offset iVal
mov [ebx], ax
mov [ebx+2], dx
```

- iVal = iVal1 x bVal

```
mov al, bVal
cbw      ; ax = bVal
cwd      ; eax = bVal
imul iVal1 ; edx:eax
mov iVal, eax ; edx is lost
```

- ascint (↓ string): int
intasc (↓ string, in): void
gets
puts

extern ascint: Near32, intasc: Near32

for ascint - can only accept numbers

sets carry flag if it's an invalid string
sets overflow flag if it's too big

page 6

call

jc

; use jump carry

jo

; use jump overflow

mov iVal, eax

add esp, 4

jmp

Assembly

Page 1

- look at questions in Chapter 3, 4, & 5 will be on Exam
- no division on the Exam

$$iVal = sVal1 \times sVal2$$

```

mov ax, sVal1
imul sVal2           ; dx:ax
x mov iVal1, eax      mov word ptr Val1, ax      or
                      mov word ptr Val2, dx

```

$$qVal = bVal1 \times iVal2$$

```

mov al, bVal1
cbw                ; ax = bVal1
cwd                ; eax = bVal1
x mov ebx, dword ptr iVal2      ; cannot go down
imul ebx           ; edx:eax
mov ebx, offset qVal
mov [ebx], eax
mov [ebx+4], edx

```

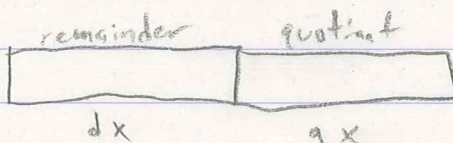
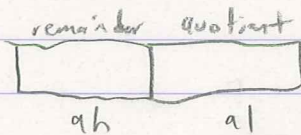
- Division - `idiv` (keeps sign)

`idiv op1`

if size of `op1` = 1, then quotient is in `ax` register

if size of `op1` = 2, then quotient should be in `dx:ax`

if size = 4, then should be in `edx:eax`



- program crashes if dividing by 0, check for this
- when multiplying, you get overflow if anything other than sign bit is there
- bT byte 7
sX word 29

mov ax, sX

idiv bT

; quotient is in al, remainder is in ah

remainder quotient	
01	04

- dividend must be twice as big as quotient
2 bytes / 1 byte = 1 byte

sY = sX / bT

mov ax, sX

idiv bT

cw

mov sY, ax

- sY = bT / sX

mov al, bT

cw

; ax = bT

cw

; dx:ax = bT

idiv sX

; dx = remainder, ax = quotient

mov sY, ax

- ; $X = sY / sX$

mov ax, sY

cwd ; dx: ax

idiv sX ; dx = remainder, ax = quotient

x mov iX, dword ptr ax

cwde

mov iX, eax

- ; $X = sY / bT$

mov ax, sY

idiv bT ; ah = remainder, al = quotient

cbw ; ax = quotient

cwde ; eax = quotient

mov iX, eax ; iX = eax

- ; $Y = sY \% bVal$

mov ax, sY ; ax = sY

idiv bT ; ah = remainder, al = quotient

mov al, ah ; move remainder into al

cbw ; ax = remainder

cwde ; eax = remainder

mov iY, eax

- sY = iVal / sVal;

x mov eax, iVal ; eax = iVal

idiv sVal ; dx = remainder, ax = quotient

mov sY, ax ; sY = ax

mov dx, word ptr iVal + 2

mov ax, word ptr iVal

OR

movsx ebx, sVal

mov sY, ax

mov eax, iVal

cdq ; edx: eax = iVal

idiv ebx

```

- ; X = sVal % iVal
  mov ax, sVal
  cwde                ; eax = sVal
  cdq                 ; edx:eax = sVal
  idiv iVal           ; eax = quotient
  mov iX, edx         ; iX = remainder

```

```

- ; iVal = sVal % bVal;
  mov ax, sVal
  idiv bVal           ; ah = remainder, al = quotient
  mov al, ah          ; al = ah
  cbw                 ; ax
  cwde                ; eax
  mov iVal, eax       ; iVal = eax

```

- Exam Review

- assembling, linking, loading
- assembler gives object file - program in machine code
- linker - adds external libraries
- loaders - uses offsets to resolve actual addresses

Hollerith code for A = 12-1